Lightweight Application-level Task Migration for Mobile Cloud Computing

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Outline

- Research background and motivation
- System design and implementation
- Performance evaluation

Background

• Mobile cloud computing:

Mobile apps or widgets connect to the Cloud

Support more complex and wider range of applications



Problems

• API lock-in → Service Provider lock-in

Client-server model: restricted form of computing



Motivation:

• Migration techniques are required to dynamically move computation between mobile nodes and cloud nodes:

Low overhead:

 Especially when using in mobile nodes where processing power and resources are very limited

• Portable:

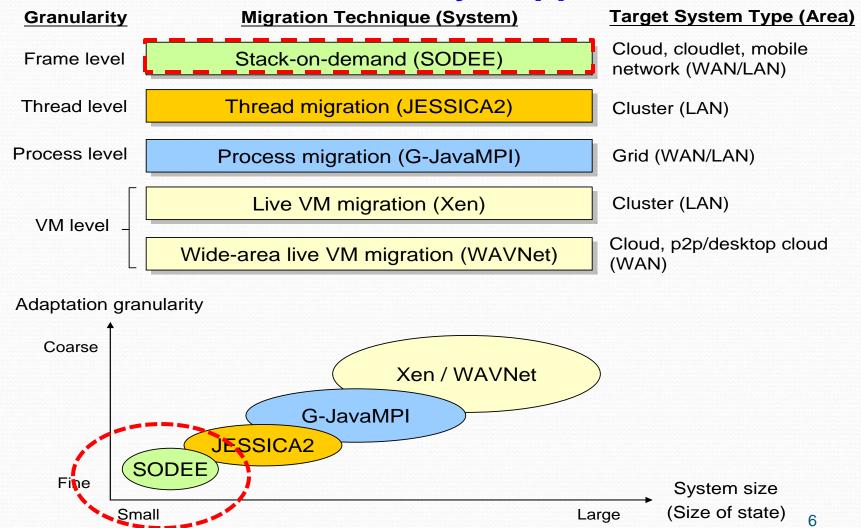
- Heterogeneous mobile nodes + Heterogeneous cloud nodes
- Task migration among mobile nodes and cloud nodes
- (Language-level virtualization for Cloud Computing)

HKU eXCloud Project Overview



(part of China National Grid (CNGrid))

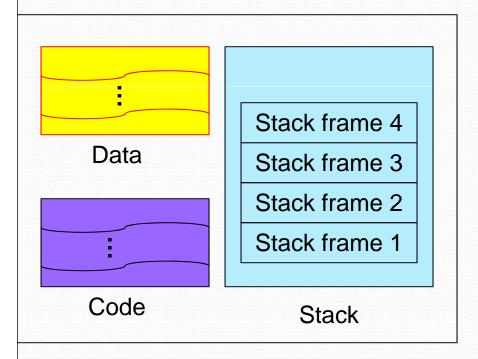
Multi-level Mobility Support

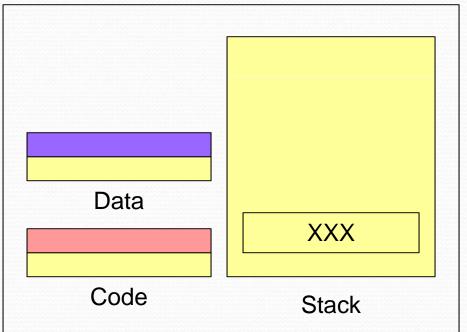


HKU eXCloud: Application Scenarios Live VM migration Stack-onover WAN (WAVNet) demand (SOD). Cloud service Method Load provider Area balancer Stack duplicate VM instances Code segments Heap for scaling Method Small Area **Thread JVM** footprint migration iOS Code **Stacks** (JESSICA) Heap **JVM** process Mobile client comm. Internet Multi-thread Java process **JVM** JVM trigger live comm. migration MPI MPI comm. maintained **Process** migration guest OS guest OS Load (G-JavaMPI) balancer Xen VM Xen VM 8 Xen-aware host OS fedoro. **Grid point Desktop PC** Overloaded 7

Stack-On-Demand (SOD): Key Ideas

Allow lightweight task migration

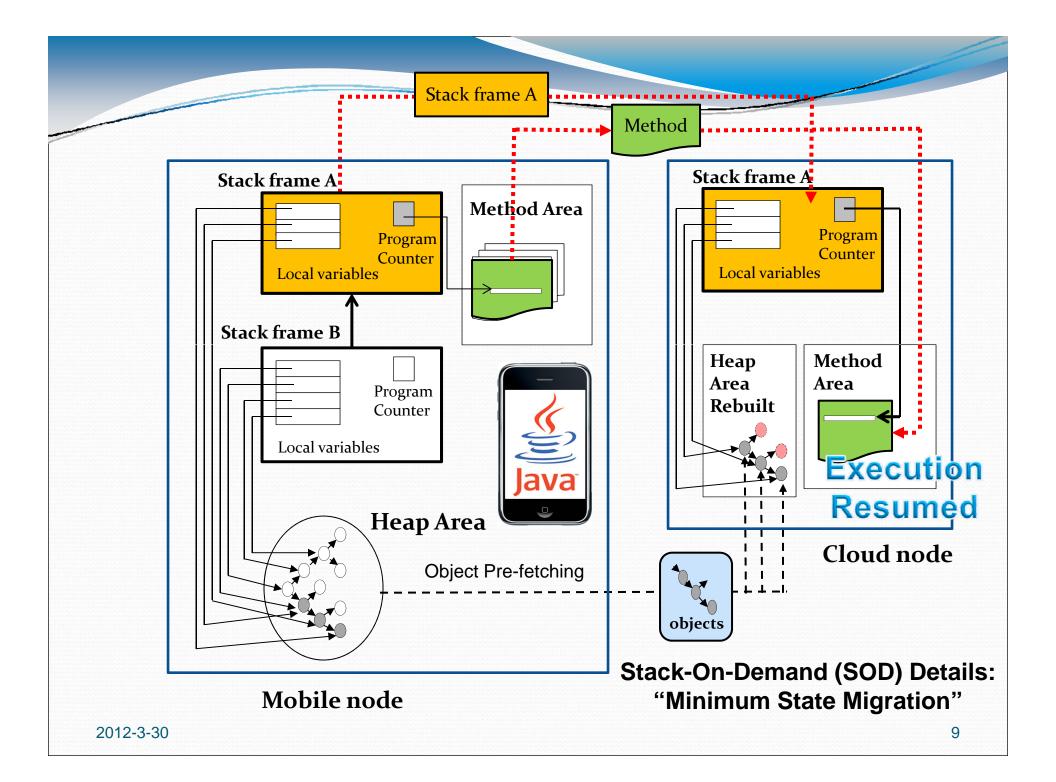




Migrating task on Source Node (a running thread)

Worker process on Destination Node

"A Stack-On-Demand Execution Model for Elastic Computing", IEEE ICPP2010.



Existing Approaches to Migrate Tasks

- At JVM level
 - Modifying JVM
 - Sumatra, ITS, CTS, JESSICA2

Requires intensive modification of JVM

Not portable for mobile nodes

- As middleware
 - through JVMTI interface
 - CIA project, G-JavaMPI
 - JVM extension
 - Mobile JikesRVM

JVMTI not available on mobile nodes

Requires certain extension of JVM

- Application-level task migration
 - Rewriting bytecode
 - JavaGoX, Brakes, JavaSplit
 - Rewriting source code
 - WASP, JavaGo



We focus on application-level task migration

Project	Level	Category	Granularity	Capturing techniques	Restoring techniques		
Merpati	JVM	Interpreter	thread	Keep state in portable format	Reconstruct based on the state		
JavaThread	JVM	Interpreter	thread	Keep state in portable format	Reconstruct based on the state		
ITS	JVM	Interpreter	thread	Keep state in portable format	Reconstruct based on the state		
CTS	JVM	JIT- compliant	thread	State in portable data structure	Reconstruct based on the state		
JESSICA2	JVM	JIT- compliant	thread	JIT recompilation	Reconstruct based on the state		
CIA	Middle- ware	extension of JVM	thread	II V VII JI + bytecode instrumentation	JVMDI + bytecode instrumentation		
Mobile JikesRVM	Middle- wa <u>re</u>	extension of JVM	thread	Use extensions of JikesRVM	Use extensions of JikesRVM		
Wasp	Application	Source-code preprocess	thread	 Java Language extended Exception (not asynchronous), but need to add migration points explicitly Part of state always saved in each migration point 	State-polling codes		
JavaGo	Application	Source-code preprocess	thread	 Java Language extended Exception (not asynchronous), but need to add migration points explicitly 	State-polling codes		
Our approach	Application	Bytecode preprocess	stack frame	Asynchronous exception (no need to add migration point and state-polling codes)	Twin Method Hierarchy (State-restoring codes executed only during restoration)		
JavaGoX	Application	Bytecode preprocess	thread	Exception (not asynchronous), but need to add migration points explicitly.	State-polling codes		
MAG/ Brakes	Application	Bytecode preprocess	thread	State-polling codes.	State-polling codes		

System Design

Design goal

- Low overhead
 - Allow lightweight task migration. Induce low overhead.
- Transparency
 - No need for users to modify their programs
- Portability
 - No need to use a specific JVM.
- Flexibility: Adaptation to new environment
 - allow to use resources in new location to utilize resources (or better resource utilization)

Common approach in application-level migration

- Use of status-polling for detecting requests
- The status-polling codes are executed even when there are no migration

Instrumentation 1: Use of status-polling for detecting requests

original statements of the function call func2()

if (isCapturing()) then

store stackframe into context store artificial PC as index value return

end if

remaining statements of the function

- 1. Status-polling codes are added for each migration point
- 2. Status-polling codes are added after each function call

- The location of inserted codes determine the migration points
- Finer granularity of migration => more insertion of status-polling codes
 => large overhead

- Use of status-polling for detecting restoration
- Status-polling codes are added at the beginning of each function call
- The status-polling codes are executed even when there are no migration

Instrumentation 2: Status-polling for detecting restoration

```
if (isRestoring()) then
  get artificial PC from context
  switch (artificial PC)
    case invokel:
      load stackframe
      goto invokel
      case ...
      ...
  end switch
end if
original statements of the function
```

Status-polling codes are added at the beginning of each function call

Our approach

- Fine-grained Task Migration
 - Among cloud nodes + Between a mobile node and a cloud node
 - Granularity : Java Stack Frame
- Two types of migration
 - Active: Triggered by migration manager
 - E.g. over loading
 - **Pro-active:** Triggered by the program itself
 - Eg. ClassNotFoundException, OutofMemoryException
 - Migration manager would then receive the requests, choose the appropriate destination and perform the migration

Task Migration

- State-capturing with Asynchronous Exception
 - Avoid status-polling (less time overheads)
 - During normal execution, as no extra codes are executed, no overhead are introduced.
 - Allow finer granularity of migration

Instrumentation with use of asynchronous exception	Capturing codes are
1. try	inserted as
2. original statements of function	exception handler
3. eatch MigrationException	
4. capture state	No significant
5. throw MigrationException	overhead introduced
6. end try	during normal

- Issues working with asynchronous exception
 - Data inconsistency
 - intermediate results are stored in operand stacks, or in native methods
 - Solution: bytecode rearrangement
 - Extra local variables are used to save intermediate results
 - Extra flags are used to inhibit migration at certain points

Deadlock

 Can lead to deadlock if asynchronous exception is used in uncontrolled manner

State Restoring with Twin Method Hierarchy

- A bytecode instrumentation technique, minimize the overhead in normal execution
- Twin Method Hierarchy
 - Keep both instrumented and original methods
 - Normal execution: original methods
 - Restoration: the instrumented methods with restoration statements are executed.
 - Checking statements are added at the beginning of the duplicated functions
 - When restoration is completed, the original method will be executed

Example

```
void func1(){
  func2();
  return;
}
```



- 1. During normal execution, original method func1() and func2() are executed => no overhead introducted
- 2. After restoration, method func1() and func2() are executed => no overhead introduced

```
void func1()
                  Original method
 func2();
 return;
                        Method used during
void SOD func1()){
                        restoration
 if (isRestoring()) {
    restore state();
     if (need restore other frame)
        goto Label1
    else
        goto previously suspended location
             original method is executed after
 func2();
             restoration has been done
Labe 12:
 return;
label1:
 SOD_func2();
                Instrumented method is executed
 goto Label2
                during restoration only
```

Performance Evaluation

Cloud server nodes

- Each node: 2 x Intel E5540 4-Core Xeon 2.53 GHz CPU, 32GB DDR3 RAM,
- OS: Fedora 11 x86_64
- JVM: Sun JDK 1.6 (64 bit)
- Network: Gigabit Ethernet

Mobile nodes

- **iPhone 4 handset**: 800MHz CPU, 512 MB RAM
- JVM: JamVM 1.5.1b2-3, slightly modified to expose the asynchronous exception API
- Java class library: GNU Classpath 0.96.1-3
- Connected to Cluster through Wi-Fi (bandwidth controlled by a router)





Performance Evaluation

Focus on performance of task migration with SOD migration

Evaluations	Description
A	Overhead analysis in cloud nodes (No Migration)
В	Overhead analysis in mobile nodes (No Migration)
С	Migration from mobile node to cloud node (by active migration)
D	Migration from mobile node to cloud node (by pro-active migration)

Evaluation A & B: Overhead Analysis

Testing programs

App	Description	Max. stack height	Total field size (byte)
Fib	Calculate 46th Fib. No.	46	< 10
NQ	Solve N-Queens problem with board size 14	16	< 10
FFT	Calculate 256-point 2D FFT	4	> 64M

• Evaluation of Three Migration Techniques:

- SOD migration using JVMTI (SOD_JVMTI)
 - implemented as JVMTI agents
 - A middleware approach (Only available on Cloud nodes)
- SOD migration using status-checking (SOD_P)
 - implemented at application level
- SOD migration using asynchronous exception (SOD_AE)
 - implemented at application level

Evaluation A: Execution time on cloud nodes (overhead when NO migration)

	Orig	SOD_JVMTI		SOD_AE		SOD_P	
	time (s)	time (s)	overhead (%)	time (s)	overhead (%)	time (s)	overhead (%)
Fib	12.11	12.13	0.17	12.14	0.25	18.4	51.78
NQ	6.35	6.4	0.79	6.7	5.51	7.24	14.02
FFT	10.53	10.63	0.95	10.82	2.75	10.6	0.47

- SOD_JVMTI imposes the smallest overhead
 - the lower layer implementations.
 - Mobile devices do not support JVMTI
- SOD_AE is slightly higher than SOD_JVMTI (< 5%)



• Evaluation B: Execution time on mobile nodes (overhead when NO migration)

	Orig	SOE	_AE	SOD_P		
	time (s)	time (s)	overhead (%)	time (s)	overhead (%)	
Fib	10.85	10.86	0.09	15.58	43.59	
NQ	32.13	32.23	0.31	33	2.71	
FFT	5.39	5.4	0.19	5.41	0.37	



- SOD_JVMTI not reported (as JVMTI is not available for JVM in mobile devices)
- SOD_AE has the smallest overhead (<0.31%)

Evaluation C: Migration from mobile device to cloud node

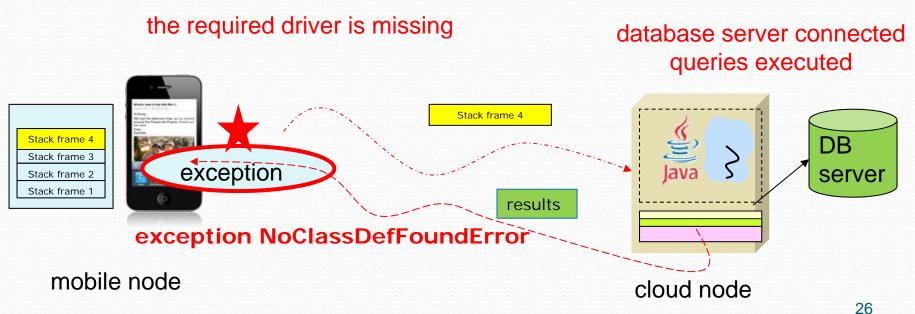
- Active migration for performance improvement
 - Migrate computation-intensive tasks from mobile devices to Cloud nodes. Upon finish of tasks, execution with data are migrated back to mobile devices.
- Performance gain: FFT: x3.8, NQ: x30, Fib: x57 times.

	exec. time w/o mig. (s)	exec. time w/ mig. (s)	Speed up	(A) capture time (ms)	(B) transfer time (ms)	(C) restore time (ms)	Total migration latency (s)
Fib	56.79	0.99	57	140.33	94.33	11.67	0.246
NQ	32.67	1.04	30	183.26	86.31	10.52	0.280
FFT	6.06	1.26	3.8	156.48	232.46	14.58	0.403

Evaluation D: Migration from mobile nodes to cloud node (Pro-Active)

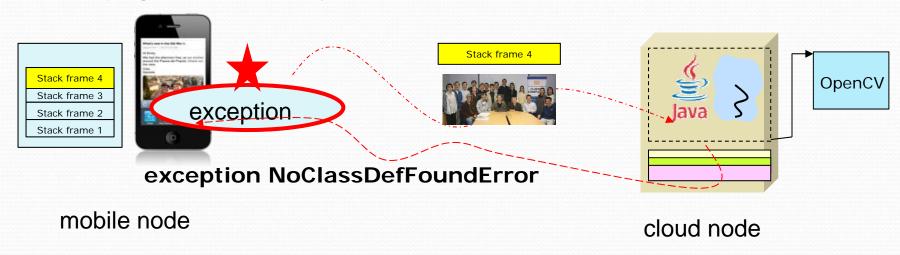
- Two applications executed in mobile nodes
 - DBRetrieve and FaceDetect
 - Both require special resources not available in mobile nodes
- **DBRetrieve** : During execution

trying to execute JDBC driver



- FaceDetect: Finds regions of faces in photos that are stored in iPhone
- Requires OpenCV library
 - open-source library for real-time computer vision and image processing
 - platform-dependent, not available in iPhone

trying to call the library OpenCV ...



Performance Evaluation



apps	capture time (ms)	transfer time (ms)	restore time (ms)	total migration latency (ms)	
DBRetrieve	85	76	6	167	
FaceDetect	103	155	7	265	

• If no SOD migration, the applications cannot be executed in mobile devices at all due to the missing resources

Conclusion and Future Work

- Java bytecode transformation technique
 - Transparent task migration in a portable and efficient manner
- Application level → Higher portability
 - Migration can take place among mobile nodes and cloud nodes
 - Does not impose significant overhead on mobile devices
- SOD Support More Flexible Computing in Cloud
 - RMI-style, process roaming, workflow model
 - Improved resource utilization
 - Avoid (API & provider) lock-in
- Future Work:
 - Killer applications of SOD?
 - Migration policies? (Resource-driven, Cost-driven, Energy-aware,..)
 - Object pre-fetching, frame-based task scheduling, ...

Thank you!

Q&A