

Mining Temporal Association Patterns under a Similarity Constraint

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Similarity-profiled Temporal Association

- A subset of items whose prevalence variation over time is similar to a reference sequence

Input

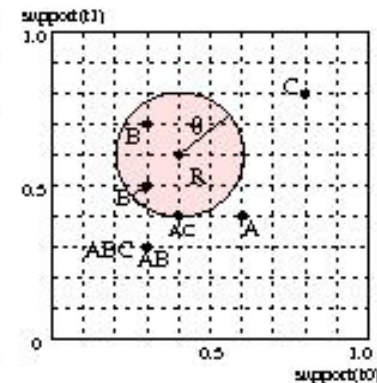
time		items	
t1	A	t2	B, C
t1	A, B, C	t2	B
t1	A, C	t2	A, B, C
t1	A	t2	A, B, C
t1	A, B, C	t2	C
t1	C	t2	A, B, C
t1	C	t2	A, C
t1	A, B, C	t2	C
t1	C	t2	B
t1	C	t2	B, C

Subject specification

Reference sequence R : $\langle 0.4, 0.6 \rangle$
 Similarity function : Euclidean distance
 Dissimilarity threshold : θ

Similarity-profiled Temporal Association Mining

itemsets	Prevalence time seq
	$\langle \text{sup}(t1), \text{sup}(t2) \rangle$
[A]	$\langle 0.6, 0.4 \rangle$
[B]	$\langle 0.3, 0.7 \rangle$
[C]	$\langle 0.8, 0.8 \rangle$
[A,B]	$\langle 0.3, 0.3 \rangle$
[A,C]	$\langle 0.4, 0.4 \rangle$
[B,C]	$\langle 0.3, 0.5 \rangle$
[A,B,C]	$\langle 0.3, 0.3 \rangle$



Output

[B] : $\langle 0.3, 0.7 \rangle (0.14)$
 [A,C] : $\langle 0.4, 0.4 \rangle (0.20)$
 [B,C] : $\langle 0.3, 0.5 \rangle (0.14)$

Motivation Examples

■ Weather-to-Sales

- ❑ Correlation between daily temperatures and merchandise sales – [Walt Disney World](#) [[NOAAEconomics](#)]
- ❑ Popular sale items during hurricane in a region – [Wal-Mart](#) [[FORTUNE Magazine](#)]
 - Flashlights, generators and tarps with bottled water
 - Strawberry Pop-Tarts with bottled water

■ Weather-to-Web Sites

- ❑ Web sites depending on weather – [[Weather.com](#)]

Motivation Examples

■ Scientific Phenomena-to-Climates

- Climate events correlated with El Nino
 - Low precipitation and low atmospheric carbon dioxide in Australia

■ Scientific Phenomena-to-Agriculture

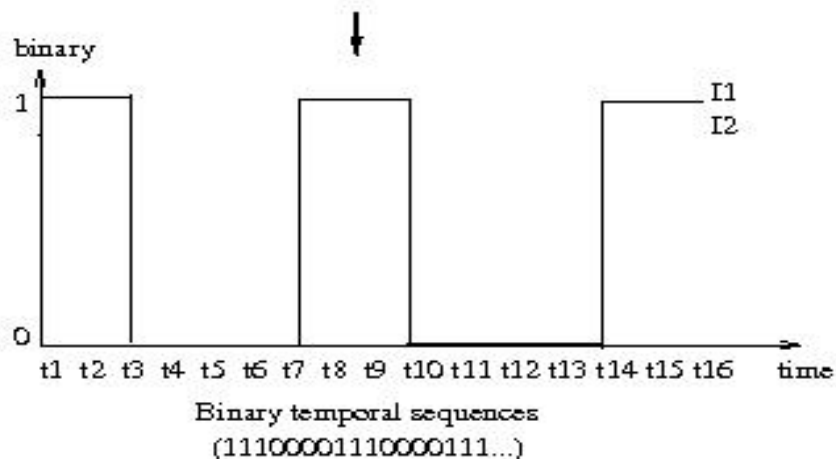
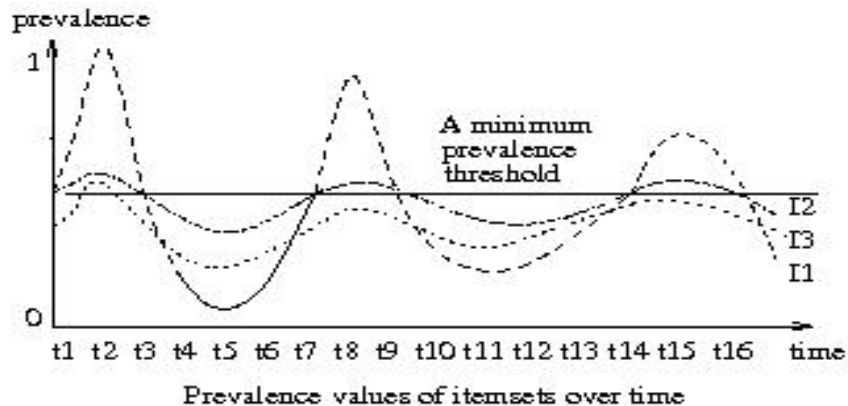
- Agricultural products under the effect of El Nino
 - Wheat and other products in Australia

Related Work

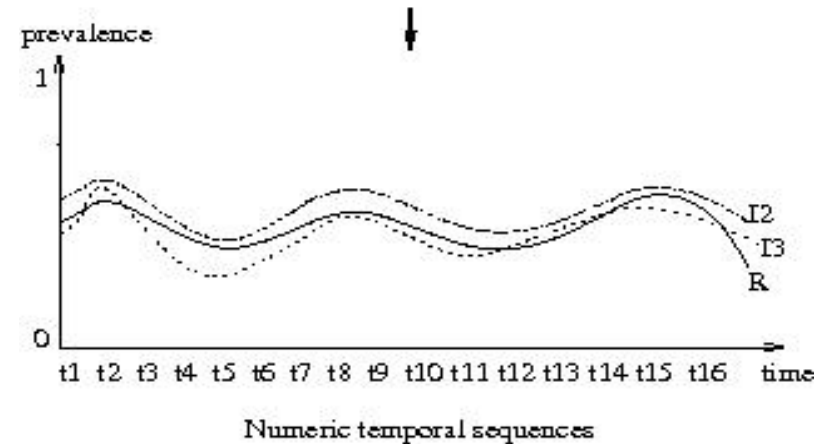
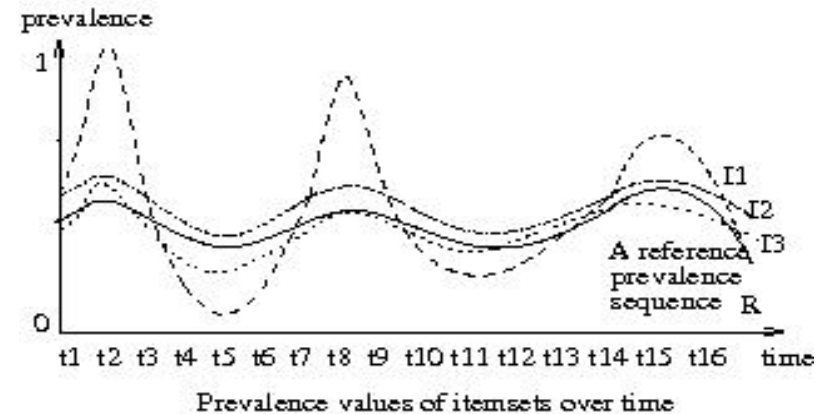
- Cyclic associations [Ozden'98, Ramaswamy'94]
 - Periodically repetitive patterns for frequent itemsets
 - “Beer” and “chips” are sold together primarily between 6PM and 9PM
- Calendar based associations [Li'03]
 - Frequent itemsets on “15th day of a March”, (*,3,15)
- User-defined temporal regulation patterns [Li'06, Bettini'98]
 - Frequent events “within 2 days” after “a rise of IBM stock”

Comparison with Related Work

■ Regulation patterns



■ Similarity patterns



Contributions

- Formulate similarity-profiled association patterns
 - User-defined temporal similarity patterns using a subset specification (i.e., a reference sequence, a similarity function, and a dissimilarity threshold)
- Explore interesting properties for efficiently mining similarity-profiled associations
- Develop the mining algorithm.
- Experimentally evaluate it with synthetic and real data sets.

Problem Definition

■ Given

- A timestamped transaction database $\mathbf{D} = \mathbf{D}_1 \cup \dots \cup \mathbf{D}_n$
 - \mathbf{D}_i is a set of transactions included in time slot i
 - Each transaction $\mathbf{d} \in \mathbf{D}$ is a tuple $\langle \text{timestamp}, \text{items} \rangle$
- A subset specification
 - A reference time sequence $\mathbf{R} = \langle r_1, \dots, r_n \rangle$
 - A similarity function $\mathbf{F}_{\text{similarity}}(\mathbf{S}_I, \mathbf{R})$, where \mathbf{S}_I is a support time sequence of itemset I
 - A dissimilarity threshold θ

■ **Find:** A set of itemsets which satisfy the given subset specification, i.e., $\mathbf{F}_{\text{similarity}}(\mathbf{S}_I, \mathbf{R}) \leq \theta$

■ **Objective:** A complete and correct result set while reducing the computation cost.

Background: Interest Measure

■ Support

- The support of itemset **I** in transaction dataset **D** is
 $support(I, D) = |\{d \in D, I \subseteq d\}| / |\{D\}|$

Transaction database	tno	items
	1	A
	2	A, B, C
	3	A, C
	4	A
	5	A, B, C
	6	C
	7	C
	8	A, B, C
	9	C
	10	C

e.g., $support(\{A\}, D) = 6/10 = 0.6$

Composite Interest Measure

- The support time sequence of itemset **I** in $\mathbf{D} = \mathbf{D}_1 \cup \dots \cup \mathbf{D}_n$
 - $\mathbf{S}_I = \langle \text{support}(I, \mathbf{D}_1), \dots, \text{support}(I, \mathbf{D}_n) \rangle$
- Dissimilarity distance between a support sequence \mathbf{S}_I and a reference sequence **R**
 - L_p norm ($p=1, 2, \dots, \infty$) based distance, e.g.,
 - L_2 norm (Euclidean distance)
 - $D(\mathbf{R}, \mathbf{S}_I) = (\sum_{t=1..n} |r_t - s_t|^2)^{\frac{1}{2}}$
 - Normalized L2 norm
 - $D(\mathbf{R}, \mathbf{S}_I) = ((\sum_{t=1..n} |r_t - s_t|^2) / n)^{\frac{1}{2}}$

Outline

- Introduction
- Problem Definition
- Related Work
- ☞ **Algorithmic Design Concept**
- Algorithm
- Experimental Results
- Conclusion

Computational Challenge

■ Naïve Approach

□ Two separate phrases

- Compute the support values of all possible itemsets at each time point, and generate their prevalence sequences
- Compare the support sequences with a reference sequence, and find similar itemsets.

□ Computationally expensive

- Exponential number of itemsets with number of item types, $2^{|n|-1}$

Our Questions

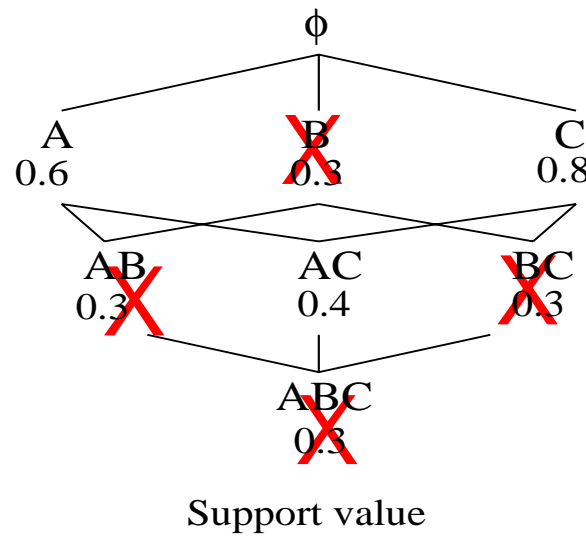
- Can we reduce the search space for only interesting patterns?
- How can we estimate the similarity distance of an itemset without the generation of the support sequence?

Background: Frequent Itemset Pruning

- Using the monotonicity of support
 - Support is **monotonically non-increasing** with the size of itemset,
i.e., $J \subseteq I$, then $\text{support}(J, \mathbf{D}) \geq \text{support}(I, \mathbf{D})$

Transaction database

tno	items
1	A
2	A, B, C
3	A, C
4	A
5	A, B, C
6	C
7	C
8	A, B, C
9	C
10	C



Frequent threshold: 0.3

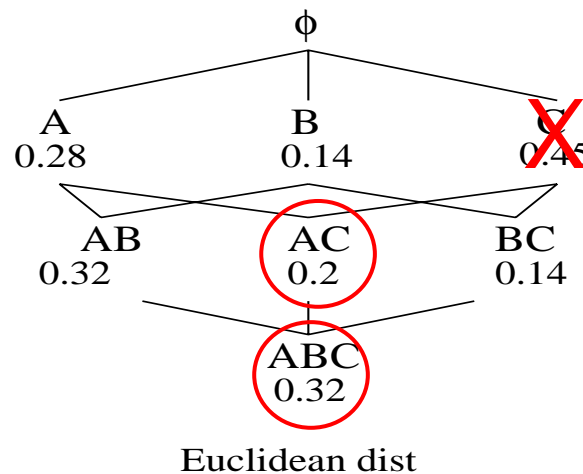
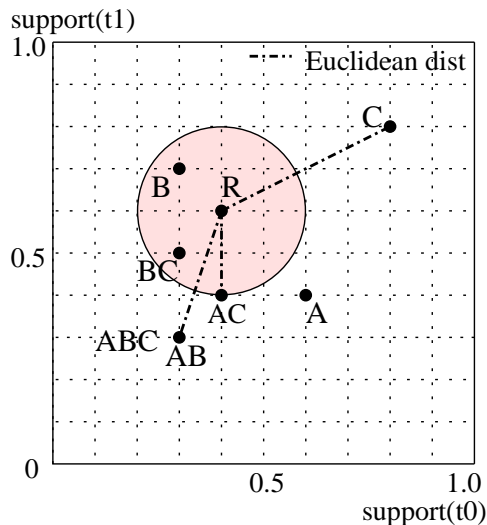
Which one is greater than 0.3?

Observation

- It is not easy to reduce our search space.

- L_p norm based distance does not show any monotonic.

e.g., $D(R, S_{\{ABC\}}) > D(R, S_{\{AC\}})$ but $D(R, S_{\{AC\}}) < D(R, S_{\{C\}})$



Let dissimilarity threshold: 0.3
Can we prune a super set of C ??

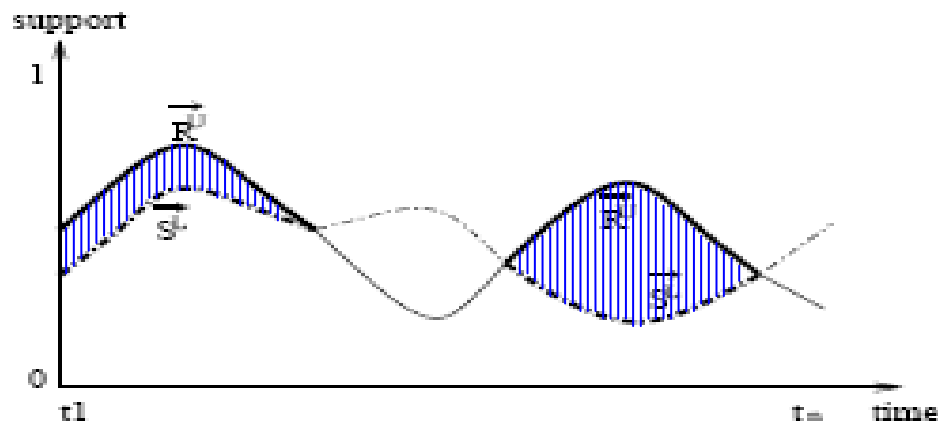
Our Approach:

Upper Lower Bounding Distance

- Let

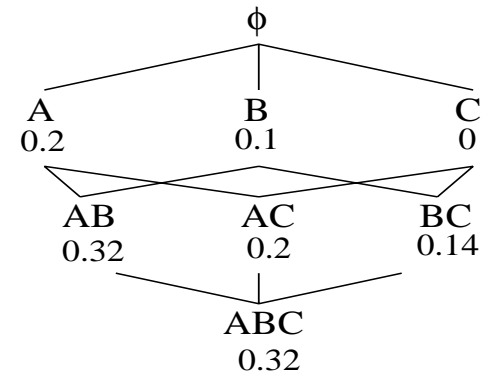
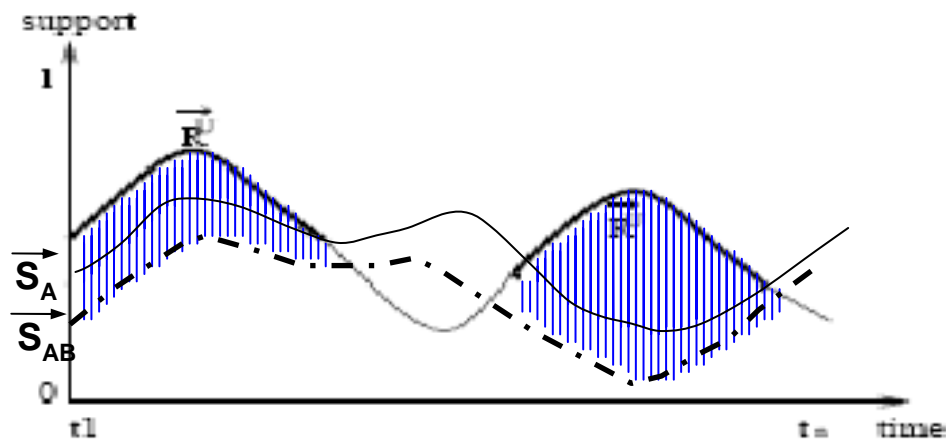
- $\mathbf{R}^U = \langle \mathbf{r}_1, \dots, \mathbf{r}_k \rangle$ be a subsequence of a reference sequence \mathbf{R} and $\mathbf{S}^L = \langle \mathbf{s}_1, \dots, \mathbf{s}_k \rangle$ be a subsequence of a support sequence \mathbf{S} , where $\mathbf{r}_i >$

- The upper lower-bounding distance between \mathbf{R} and \mathbf{S} , $D_{Ulb}(\mathbf{R}, \mathbf{S})$, is $D(\mathbf{R}^U, \mathbf{S}^L)$.



Monotonicity of Upper LBD

- The upper lower-bounding distance is **monotonically non-decreasing** with the size of the itemset.
- Proof: The support values of an itemset are **monotonically non-increasing** with the size of itemset at each time slot.



Upper lower bounding dist

$$\begin{aligned} \text{E.g., } D_{Ulb}(R, S_A) &\leq D_{Ulb}(R, S_{AB}), \\ D_{Ulb}(R, S_B) &\leq D_{Ulb}(R, S_{AB}), \end{aligned}$$

Prune by Upper LBD

- Let Itemsets $J \subseteq I$
- If $D_{Ulb}(\mathbf{R}, \mathbf{S}_J) > \theta$, always $D_{Ulb}(\mathbf{R}, \mathbf{S}_I) > \theta$
- Prune all superset of J

Our Questions

- Can we reduce the search space for only interesting patterns?
- ☞ How can we estimate the similarity distance of an itemset without the generation of the support sequence?

Upper Bound of Support Sequence

- Let

- $\mathbf{D} = \mathbf{D}_1 \cup \dots \cup \mathbf{D}_n$ be a set of disjoint transactions.
- $\mathbf{J} = \{\mathbf{J}_1, \dots, \mathbf{J}_k\}$ be a set of all size $k-1$ subsets of a size k itemset \mathbf{I} .

- Upper bound support time sequence of itemset \mathbf{I} , $\mathbf{U}_1 = \langle \mathbf{u}_1, \dots, \mathbf{u}_n \rangle$ is defined as

- $\mathbf{u}_1 = \min \{ \text{support}(\mathbf{J}_1, \mathbf{D}_1), \dots, \text{support}(\mathbf{J}_k, \mathbf{D}_1) \}$
- $\mathbf{u}_n = \min \{ \text{support}(\mathbf{J}_1, \mathbf{D}_n), \dots, \text{support}(\mathbf{J}_k, \mathbf{D}_n) \}$

itemsets	Prevalence time seq
	$\langle \text{sup}(t1), \text{sup}(t2) \rangle$
{A}	$\langle 0.6, 0.4 \rangle$
{B}	$\langle 0.3, 0.7 \rangle$
{C}	$\langle 0.8, 0.8 \rangle$
{A,B}	$\langle 0.3, 0.3 \rangle$
{A,C}	$\langle 0.4, 0.4 \rangle$
{B,C}	$\langle 0.3, 0.5 \rangle$
{A,B,C}	?

- E.g., $\mathbf{U}_{ABC} = \langle \mathbf{u}_1, \mathbf{u}_2 \rangle = \langle 0.3, 0.3 \rangle$
 $\mathbf{u}_1 = \min \{ \text{supp}(\mathbf{AB}, \mathbf{D}_1), \text{supp}(\mathbf{AC}, \mathbf{D}_1), \text{supp}(\mathbf{BC}, \mathbf{D}_1) \}$
 $\mathbf{u}_2 = \min \{ \text{supp}(\mathbf{AB}, \mathbf{D}_2), \text{supp}(\mathbf{AC}, \mathbf{D}_2), \text{supp}(\mathbf{BC}, \mathbf{D}_2) \}$

Upper Bound of Support Sequence

- Let
 - $\mathbf{D}=\mathbf{D}_1 \cup \dots \cup \mathbf{D}_n$ be a set of disjoint transactions.
 - $\mathbf{J}=\{\mathbf{J}_1, \dots, \mathbf{J}_k\}$ be a set of all size $k-1$ subsets of a size k itemset \mathbf{I} .

- Lower bound support time sequence of itemset \mathbf{I} , $\mathbf{L}_1 = \langle \mathbf{I}_1, \dots, \mathbf{I}_n \rangle$ is defined as

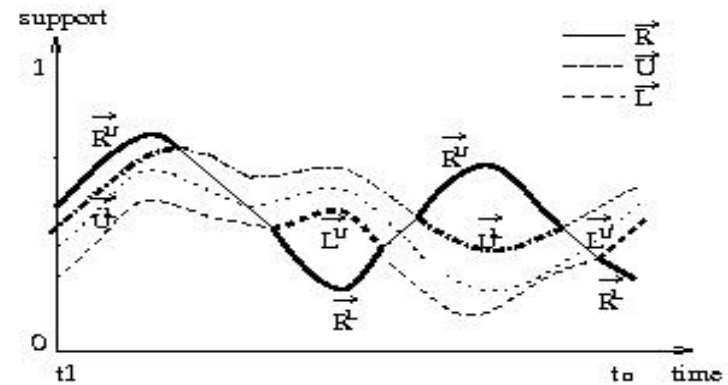
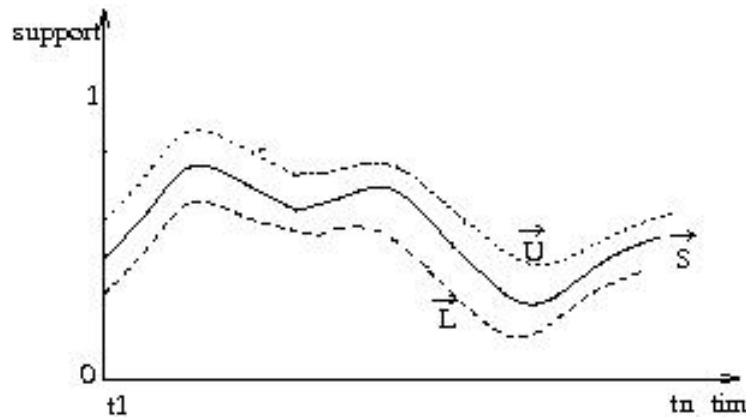
- $\mathbf{I}_1 = \max \{(\text{support}(\mathbf{J}_1, \mathbf{D}_1) + \text{support}(\mathbf{I}-\mathbf{J}_1, \mathbf{D}_1) - 1), \dots, (\text{support}(\mathbf{J}_k, \mathbf{D}_1) + \text{support}(\mathbf{I}-\mathbf{J}_k, \mathbf{D}_1) - 1), 0\}$
- $\mathbf{I}_n = \max \{(\text{support}(\mathbf{J}_1, \mathbf{D}_n) + \text{support}(\mathbf{I}-\mathbf{J}_1, \mathbf{D}_n) - 1), \dots, (\text{support}(\mathbf{J}_k, \mathbf{D}_n) + \text{support}(\mathbf{I}-\mathbf{J}_k, \mathbf{D}_n) - 1), 0\}$

itemsets	Prevalence time seq
	$\langle \text{sup}(t1), \text{sup}(t2) \rangle$
{A}	$\langle 0.6, 0.4 \rangle$
{B}	$\langle 0.3, 0.7 \rangle$
{C}	$\langle 0.8, 0.8 \rangle$
{A,B}	$\langle 0.3, 0.3 \rangle$
{A,C}	$\langle 0.4, 0.4 \rangle$
{B,C}	$\langle 0.3, 0.5 \rangle$
{A,B,C}	?

- E.g., $\mathbf{L}_{\text{ABC}} = \langle \mathbf{I}_1, \mathbf{I}_2 \rangle = \langle 0.1, 0.1 \rangle$

$$\mathbf{u}_1 = \max \{(\text{supp}(\mathbf{AB}, \mathbf{D}_1) + \text{supp}(\mathbf{C}, \mathbf{D}_1) - 1), (\text{supp}(\mathbf{AC}, \mathbf{D}_1) + \text{supp}(\mathbf{B}, \mathbf{D}_1) - 1), (\text{supp}(\mathbf{BC}, \mathbf{D}_1) + \text{supp}(\mathbf{A}, \mathbf{D}_1) - 1), 0\}$$

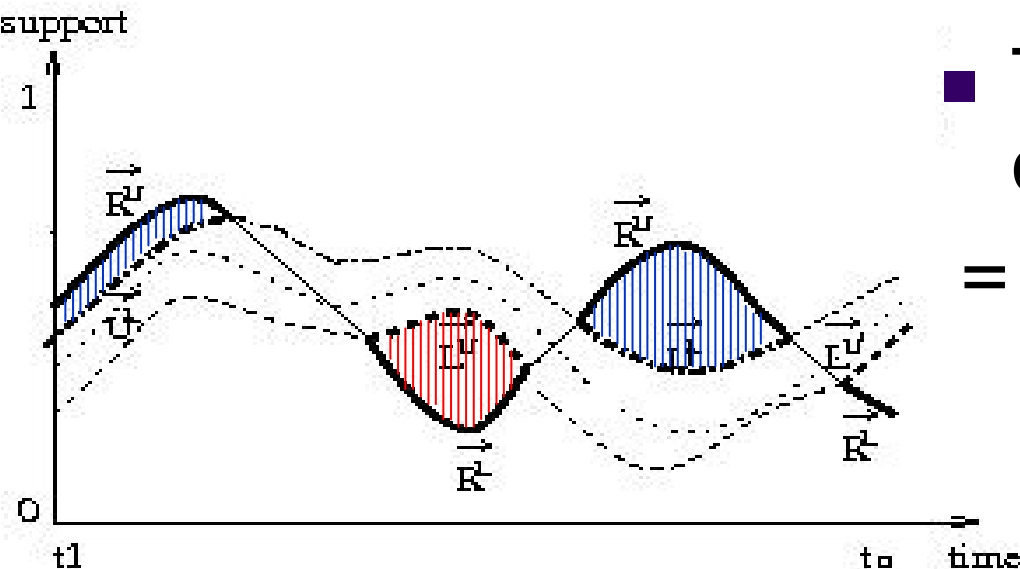
Subsequences for Lower Bounding Distance



- $\vec{R}^U = \langle r_1, \dots, r_k \rangle$ be a subsequence of \vec{R} and $\vec{U}^L = \langle u_1, \dots, u_k \rangle$ be a subsequence of \vec{U} , where $r_i > u_i$
- $\vec{R}^L = \langle r_1, \dots, r_k \rangle$ be a subsequence of \vec{R} and $\vec{L}^U = \langle l_1, \dots, l_k \rangle$ be a subsequence of \vec{L} , where $r_i > l_i$

Lower Bounding Distance

- The **upper lower-bounding distance** between R and U , $D_{Ulb}(R, U)$ is defined to $D(R^U, U^L)$.
- The **lower lower-bounding distance** between R and L , $D_{Llb}(R, L)$ defined to $D(R^L, L^U)$.



- The **lower-bounding distance**, $D_{lb}(R, U, L) = D_{Ulb}(R, U) + D_{Llb}(R, L)$

Prune by Lower Bounding Distance

- The lower bounding distance $D_{lb}(\mathbf{R}, U_I, \mathbf{L}_I)$ is always not greater than true distance $D(\mathbf{R}, S_I)$.
- So, If $D_{lb}(\mathbf{R}, U_I, \mathbf{L}_I) > \theta$, $D(\mathbf{R}, S_I) > \theta$
- Prune itemset I

Database Scan Strategy

- Lattice-dominant scan
- Snapshot-dominant scan

Time-stamped transaction database

time	items	time	items
t0	A	t1	B, C
t0	A, B, C	t1	B
t0	A, C	t1	A, B, C
t0	A	t1	A, B, C
t0	A, B, C	t1	C
t0	C	t1	A, B, C
t0	C	t1	A, C
t0	A, B, C	t1	C
t0	C	t1	B
t0	C	t1	B, C

Time-stamped transaction database

time	items	time	items
t0	A	t1	B, C
t0	A, B, C	t1	B
t0	A, C	t1	A, B, C
t0	A	t1	A, B, C
t0	A, B, C	t1	C
t0	C	t1	A, B, C
t0	C	t1	A, C
t0	A, B, C	t1	C
t0	C	t1	B
t0	C	t1	B, C

Outline

- Introduction
- Problem Definition
- Related Work
- Algorithmic Design Concept
- ☞ **Algorithm**
- Experimental Results
- Conclusion

Similarity-Profiled temporal Association MINing methods


- Two algorithms by different database scan methods
 - L-SPAMINE: Lattice-dominant SPAMINE
 - S-SPAMINE: Snapshot-dominant SPAMINE

Algorithm (L-SPAMINE)

■ Input

- A time-stamped dataset
- A reference sequence, A similarity function, and A threshold

■ Procedure

- Generate size K candidate itemsets
 - Prune if **any subset's** $D_{Ulb} < \text{threshold}$
 - Estimate **upper and lower bound sequences** of candidates
 - Filter candidates using $D_{lb} (=D_{Ulb} + D_{Lib})$
 - Scan database and generate true support sequences
 - Find similar itemsets having $D < \text{threshold}$
 - Keep size K itemsets having $D_{Ulb} < \text{threshold}$
 - $K=K+1$
- 

L-SPAMINE Trace

Transaction database	time	items	time	items
	t1	A	t2	B, C
	t1	A, B, C	t2	B
	t1	A, C	t2	A, B, C
	t1	A	t2	A, B, C
	t1	A, B, C	t2	C
	t1	C	t2	A, B, C
	t1	C	t2	A, C
	t1	A, B, C	t2	C
	t1	C	t2	B
	t1	C	t2	B, C

		Size 1		
Reference sequence		Support sequences		
		A	B	C
t1	0.4	0.6	<u>0.3</u>	0.8
t2	0.6	<u>0.4</u>	0.7	0.8
Upper LB dist:		0.20✓	0.10✓	0 ✓
True dist:		0.28✗	0.14✓	0.45✗

Dissimilarity threshold : 0.2

Similarity Function: Euclidean

L-SPAMINE Trace

Transaction database

time	items	time	items
t1	A	t2	B, C
t1	A, B, C	t2	B
t1	A, C	t2	A, B, C
t1	A	t2	A, B, C
t1	A, B, C	t2	C
t1	C	t2	A, B, C
t1	C	t2	A, C
t1	A, B, C	t2	C
t1	C	t2	B
t1	C	t2	B, C

Dissimilarity threshold : 0.2
Similarity Function: Euclidean

Size 2

Reference sequence	Upper bound sequences		
	A B	A C	B C
t1	<div><div>0.4</div></div>	0.6	<u>0.3</u>
t2	<div><div>0.6</div></div>	<u>0.4</u>	0.7
Upper LB dist:	0.22	0.20✓	0.10✓

Lower bound sequences

A B	A C	B C
	0.4	0.1
	0.2	0.5
Lower LB dist:	0.0	0.0
LB dist:	0.20✓	0.10✓

Size 2

Support sequences	A C	B C
	0.6	<u>0.3</u>
	<u>0.4</u>	<u>0.5</u>
Upper LB dist:	0.20✓	0.14✓
True dist:	0.20✓	0.14✓

Size 3

A B C

* Similar itemsets:
 {B} : <0.3., 0.7> (0.14)
 {A,C}: <0.6., 0.4> (0.2)
 {B,C}: <0.3., 0.5> (0.14)

Experiment

■ Datasets

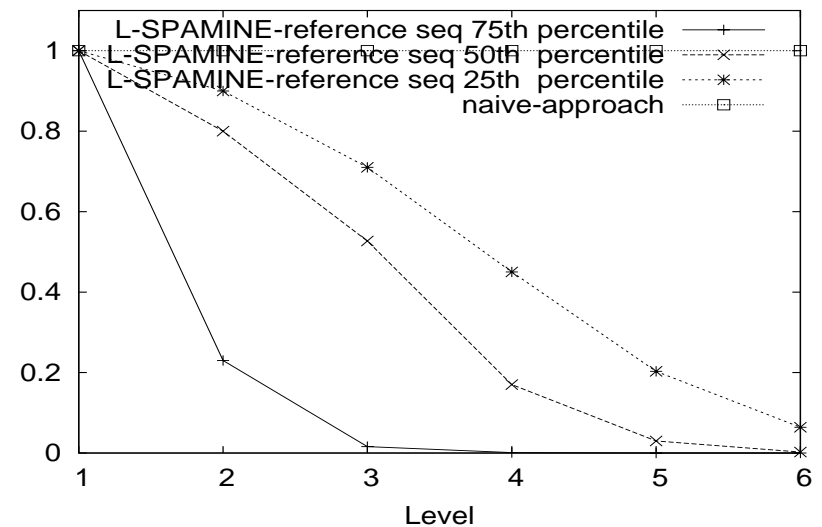
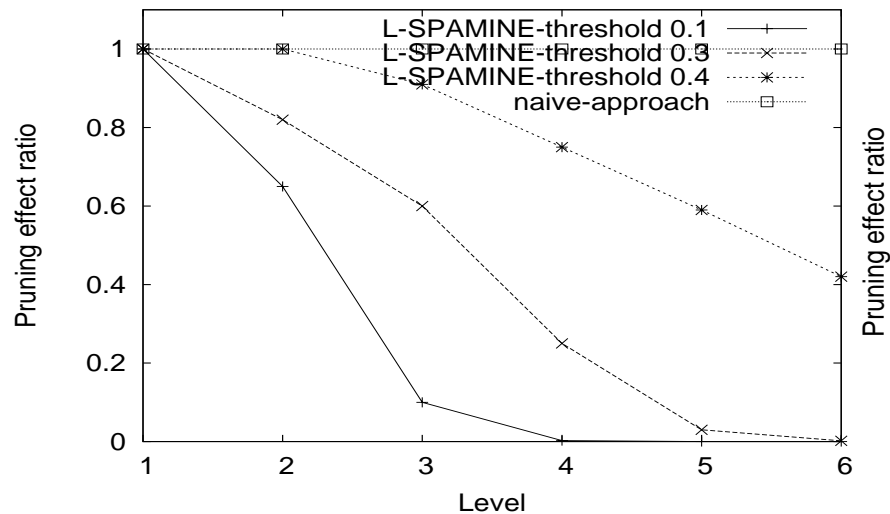
- ❑ Synthetic datasets: a modified IBM data generator
- ❑ Real dataset: Earth Climate
- ❑ Query sequences: randomly chosen in different quintiles of supports

■ Test cases

- ❑ Effect of lower bounding distance
- ❑ Effect of database scanning method
- ❑ Effect of number of items
- ❑ Effect of number of time slots
- ❑ Experiment with a real dataset

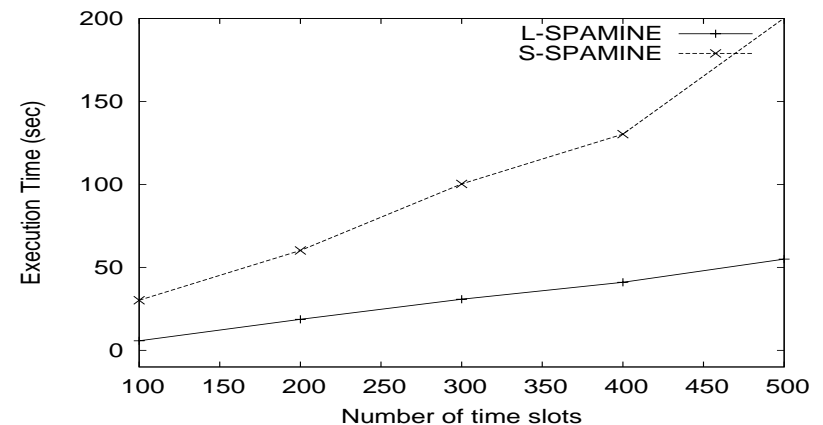
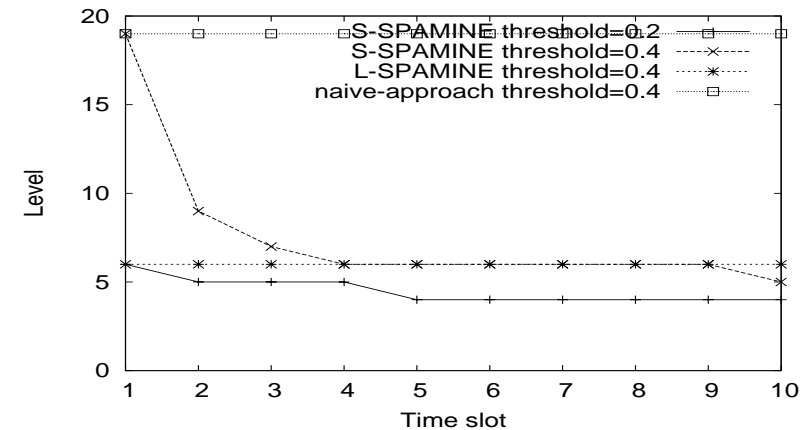
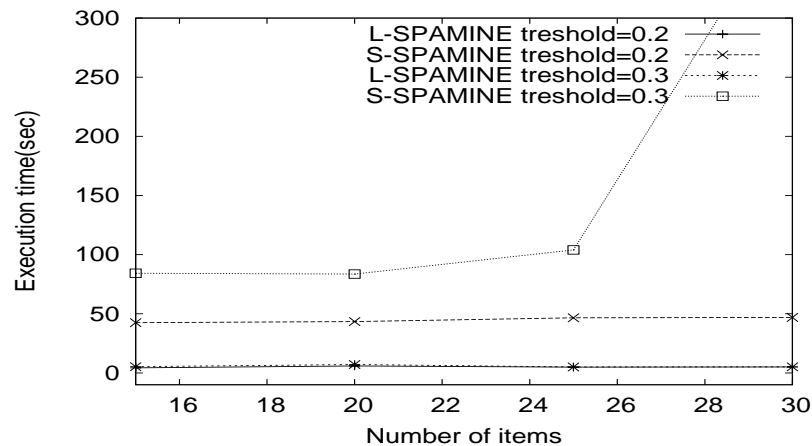
Experiment Results

- Effect of lower bounding distance pruning
 - TD100-D1-L10-I20-T100
 - Pruning effect ratio : the number of candidate itemsets which need database scan over the total number of possible itemsets per level



Experiment Results

■ Effect of different scanning (TD100-D1-L10-I20-T100)



■ Effect of number of items (TD100-D1-L10-I*-T10)

■ Effect of number of time slots (TD*-D1-L6-I20-T*)

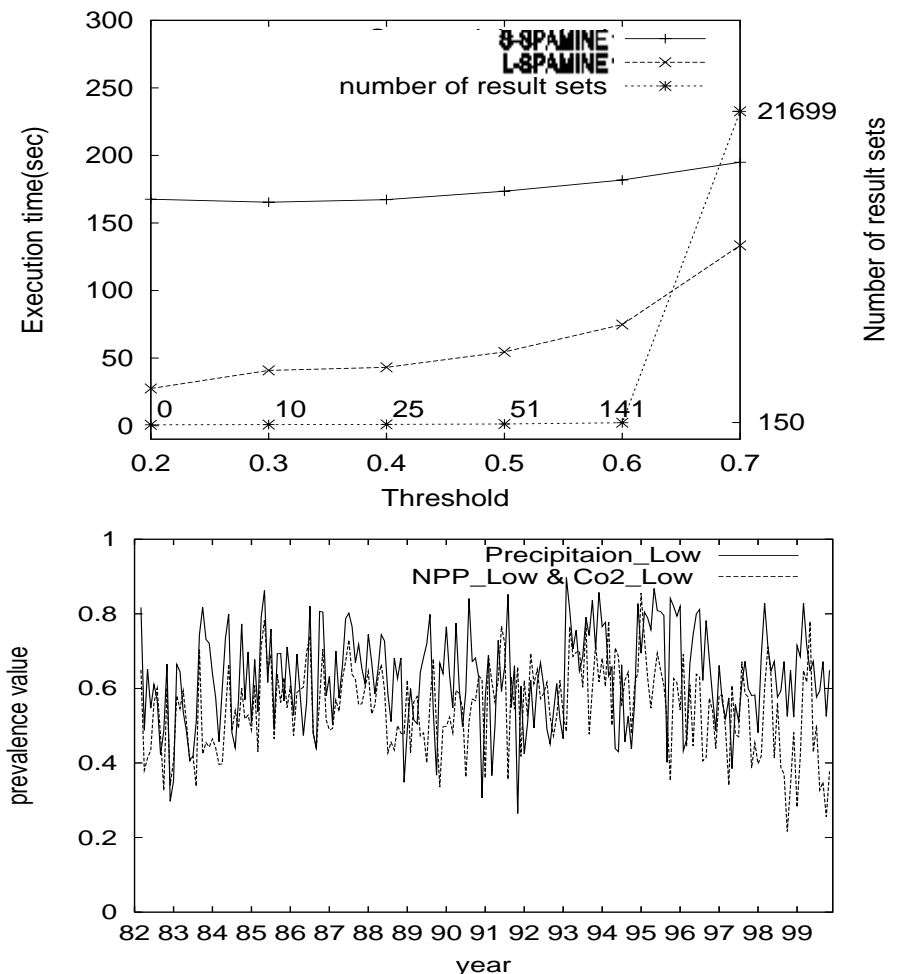
Experiment with a real dataset

■ Dataset: Earth Climate

- # of items: 50,
- # of time slots: 214
- # of transaction per time slot: 2827,
- Total # of transaction: 64,978

■ Reference sequences

- SOI index
 - Normalization to 0 to 1 range.
- Prevalence sequence of low participation



Conclusion

■ Summary

- ❑ Formulate the problem of mining similarity-profiled temporal association patterns
- ❑ Propose a novel algorithm
 - Substantially reduce the search space by pruning candidate itemsets using lower bounding distance and the monotonicity of upper lower bounding distance
- ❑ Experimentally evaluate the algorithm

■ Future Work

- ❑ Explore different similarity measures with different similarity models.